

An algebraic internal groupoid model of Martin-Löf type theory

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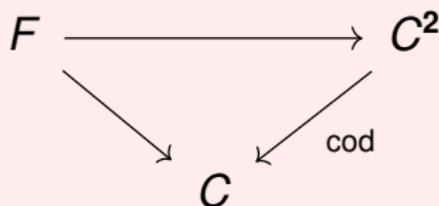


Comprehension category

Recall:

Definition ([Jac93])

A comprehension category is a strictly commutative diagram of functors



such that $F \rightarrow C$ is a Grothendieck fibration and $F \rightarrow C^2$ is a cartesian functor.

such data gives a (weak) model of MLTT.

Example: the groupoid model

An isofibration in **Gpd** is

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i.e. the map $(F_1, d_0) : A_1 \rightarrow B_1 \times_{B_0} A_0$ has a section.

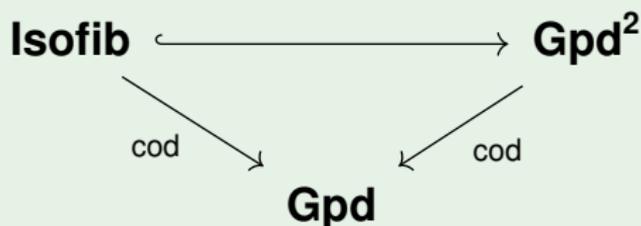
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Example ([HS98])



forms a comprehension category.

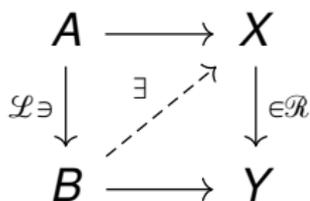
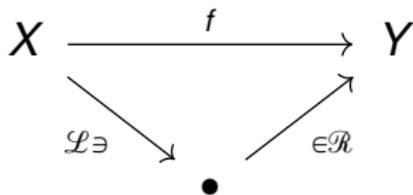
Weak Factorisation Systems

Many examples of comprehension categories arise from weak factorisation systems.

Definition

A *weak factorisation system* (wfs) on a category \mathbf{C} is a pair $(\mathcal{L}, \mathcal{R})$ of classes of morphisms in \mathbf{C} such that:

- 1 Every map $f : X \rightarrow Y$ can be factorised as a map in \mathcal{L} followed by a map in \mathcal{R} .
- 2 $\mathcal{L} = {}^{\perp} \mathcal{R}$ and $\mathcal{R} = \mathcal{L}^{\perp}$.



Examples of WFSs

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(injective-on-object equivalences, isofibrations) form a weak factorisation system on **Gpd**.

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... but can't often be equipped with coherent path objects.
A fix of this was suggested by Garner in moving to the algebraic setting.

Type theoretic awfs

Definition ([GL23])

A *type-theoretic algebraic weak factorisation system* on a category \mathbf{C} is a pair (\mathbb{L}, \mathbb{R}) of a comonad and a monad on \mathbf{C}^{\rightarrow} such that $(\overline{\mathbb{L}\text{-Coalg}}, \overline{\mathbb{R}\text{-Alg}})$ is a wfs on \mathbf{C} with some extra structure and satisfying certain conditions.

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Theorem ([GL23], Theorem 4.12)

Let (\mathbb{C}, \mathbb{F}) be a type theoretic algebraic weak factorisation system. Then the right adjoint splitting of the comprehension category associated to the awfs is equipped with strictly stable choices of Σ , Π and Id -types i.e. it forms a model of MLTT.

The \mathbb{F} -algebras model the dependent types.

Cloven isofibrations

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Definition

A **cloven** isofibration is a pair (F, s) in which the map $s : B_1 \times_{B_0} A_0 \rightarrow A_1$ is a chosen section of (F_1, d_0) .

The algebraic groupoid model

Proposition

There is a monad $\mathbb{F} : \mathbf{Gpd}^2 \rightarrow \mathbf{Gpd}^2$ such that $\mathbb{F}\text{-Alg} \cong \mathbf{ClovenIsfibrations}$.

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There is a type theoretic AWFS involving \mathbb{F} . Hence cloven isofibrations model MLTT.

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In the type theory, $(F, s) \neq (F, t)$ for $s \neq t$.

Internal groupoids

Definition

A *small groupoid* is:

$$\begin{array}{ccccc} & & \xrightarrow{p_1} & & \xrightarrow{d_1} \\ \dots & \longrightarrow & C_1 \times_{C_0} C_1 & \xrightarrow{m} & C_1 & \xleftarrow{i} & C_0 \\ & & \xrightarrow{p_2} & \uparrow & \xrightarrow{d_0} \\ & & & (-)^{-1} & \end{array}$$

Where $C_0, C_1 \in \mathbf{Set}$. These are the objects of a $(2, 1)$ -category **Gpd**.

Internal groupoids

Definition

Let \mathcal{E} be a category with pullbacks.

A **groupoid internal to \mathcal{E}** is:

$$\begin{array}{ccccc} \dots & \longrightarrow & C_1 \times_{C_0} C_1 & \begin{array}{c} \xrightarrow{p_1} \\ \xrightarrow{m} \\ \xrightarrow{p_2} \end{array} & C_1 & \begin{array}{c} \xleftarrow{d_1} \\ \xleftarrow{i} \\ \xleftarrow{d_0} \end{array} & C_0 \\ & & & & \uparrow & & \\ & & & & (-)^{-1} & & \end{array}$$

Where $C_0, C_1 \in \mathcal{E}$.

These are the objects of a $(2, 1)$ -category **Gpd**(\mathcal{E}).

Internal cloven isofibrations

Definition

An internal cloven isofibration is a pair (F, s) in which the map $s : B_1 \times_{B_0} A_0 \rightarrow A_1$ is a chosen section of $(F_1, d_0) : A_1 \rightarrow B_1 \times_{B_0} A_0$.

The internal algebraic groupoid model

Let \mathcal{E} be a locally cartesian closed lex extensive categories in which the forgetful functor $\mathcal{U} : \mathbf{Cat}(\mathcal{E}) \rightarrow \mathbf{Gph}(\mathcal{E})$ has a left adjoint:

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Theorem

There is a monad $\mathbb{F} : \mathbf{Gpd}(\mathcal{E})^2 \rightarrow \mathbf{Gpd}(\mathcal{E})^2$ such that $\mathbb{F}\text{-Alg} \cong \mathbf{ClovenIsofibrations}$. Moreover, there is a type theoretic AWFS involving \mathbb{F} . Hence internal cloven isofibrations model MLTT.

Non-examples

Non-examples: $\mathcal{C} = \mathbf{Cat}, \mathbf{Cat}(\mathcal{C}), \mathbf{Ab} \dots$

Examples

Locally cartesian closed lexensive categories in which the forgetful functor $\mathcal{U} : \mathbf{Cat}(\mathcal{E}) \rightarrow \mathbf{Gph}(\mathcal{E})$ has a left adjoint:

- **Set**

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- Arithmetic Π -pretoposes [Mai10].
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- **Asm** (cf. the effective topos [Hyl88])...

Future work

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- Moreover, it forms a model of MLTT.
- We can find a modest discrete opfibration classifier in $\mathbf{Gpd}(\mathbf{Asm})$ (cf. [Web07]). In the type theory, this gives a univalent universe of modest 0-types.
- Moreover, we show that modest discrete opfibrations form a 2-category with a class of small discrete opfibrations (cf. [JM95]).

Future work II

Can we do this for $\mathbf{s}^{\mathcal{E}} := [\Delta^{\text{op}}, \mathcal{E}]$ and/or $[\square^{\text{op}}, \mathcal{E}]$?

The algebraic internal groupoid model of Martin-Löf type theory, 2025.



<https://arxiv.org/abs/2503.17319>

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