

Representing type theories in two-level type theory

Nicolai Kraus

jww **Tom de Jong**

TYPES 2025

Glasgow, 13 June 2025

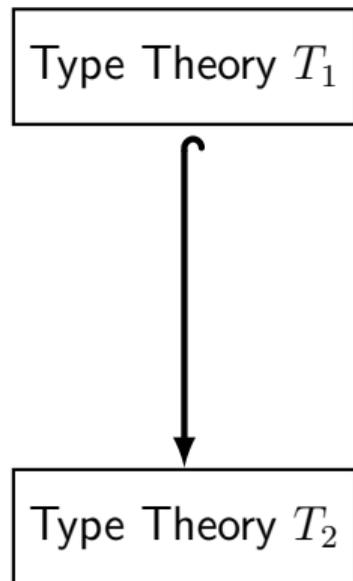
Goal: Understanding Extensions

Situation:

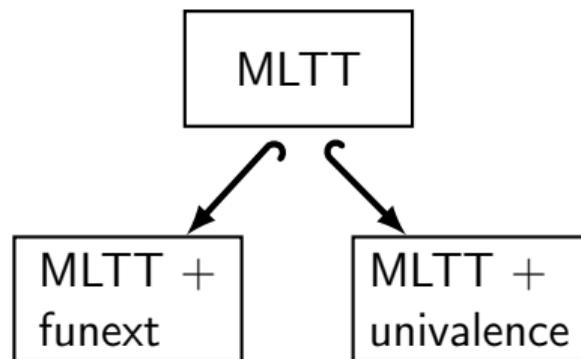
- ▶ We understand T_1 .
- ▶ We want to understand T_2 .
- ▶ We want to compare T_2 to another extension of T_1 .

Two kinds of extensions:

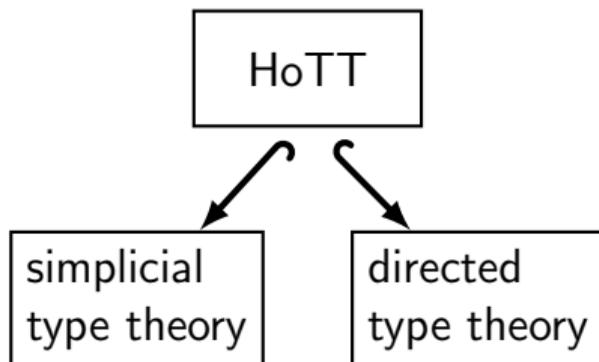
- ▶ axiomatic extensions (easier to understand, can use proof assistants for T_1)
- ▶ structural extensions (trickier)



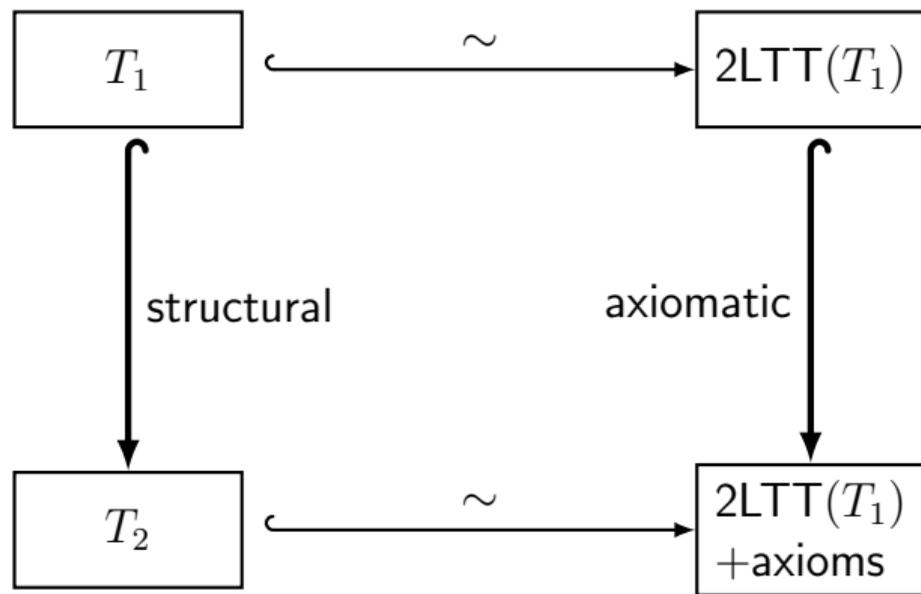
Axiomatic extensions:



Structural extensions:



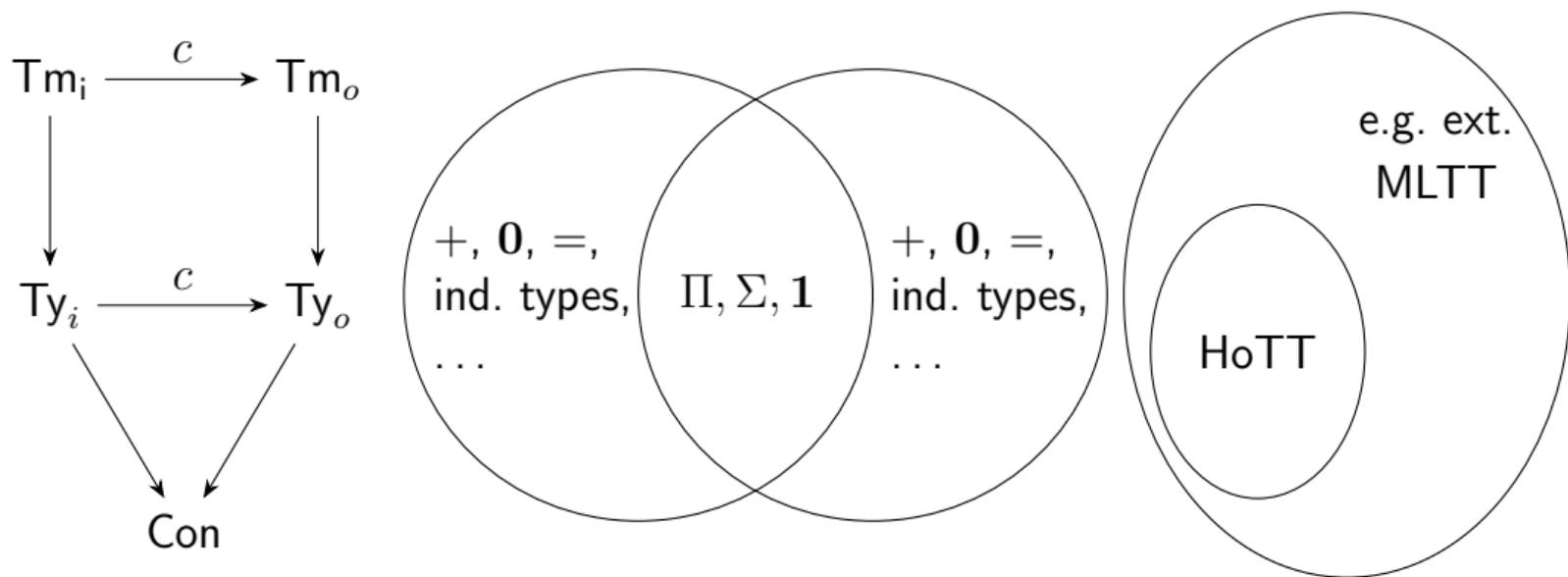
Approach: Make Extensions Axiomatic



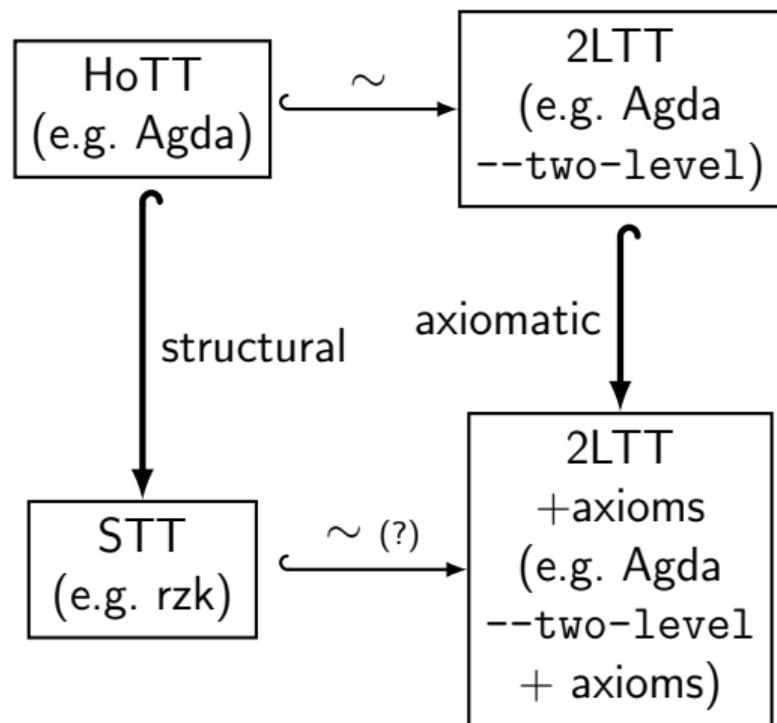
- ▶ We want to use two-level type theory (2LTT) to represent other theories.
- ▶ The vertical extensions are ideally conservative (i.e. don't change what is provable).
- ▶ This may lose computational properties.

Two-level type theory [Voe13; Cap17; ACK18; Ann+23]

- ▶ Two type theories (*inner/fibrant* and *outer/exo* theory)
- ▶ Inner theory is the type theory of interest; outer theory is “just” auxiliary language



Example: simplicial type theory in 2LTT



- ▶ Riehl and Shulman's *simplicial type theory (STT)* [RS17] is HoTT with two additional components:
 - ▶ additional context layers to talk about *shape inclusions* and *extension types*;
 - ▶ simplicial shapes.
- ▶ Remainder of this talk: model STT in 2LTT.

Example: simplicial type theory in 2LTT; extension types

In simplicial TT, assume:

- ▶ $\Phi \subset \Psi$ are shapes (defined using the new context layers)
- ▶ A is a type on the “big” shape Ψ
- ▶ a is a term of A on the “small” shape Φ

Then: $\langle \Pi_{t:\Psi} A(t) |_a^\Phi \rangle$ is the type of extensions of a .

In 2-level type theory:

- ▶ $i : L \rightarrow K$ is *cofibration* if $f\hat{\mathfrak{H}}_-$ preserves [trivial] fibrations.
- ▶ This means: For any fibrant family $Y : M \rightarrow \mathcal{U}$ and strictly commuting squares the type of (d) is fibrant (and contractible if Y is).

$$\begin{array}{ccc} L & \xrightarrow{u} & \Sigma_M Y \\ i \downarrow & \nearrow d & \text{pr} \downarrow \\ K & \xrightarrow{v} & M \end{array}$$

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- ▶ Special case:

$$\begin{array}{ccc} \Phi & \xrightarrow{a} & \Sigma_\Psi Y \\ i \downarrow & \nearrow d & \text{pr} \downarrow \\ \Psi & \xrightarrow{id} & \Psi \end{array}$$

Shape inclusions of simplicial type theory = cofibrations of 2LTT

Extension types of simplicial type theory = properties of cofibrations

Example: simplicial type theory in 2LTT; simplicial shapes

Second ingredient of simplicial type theory: a directed interval.

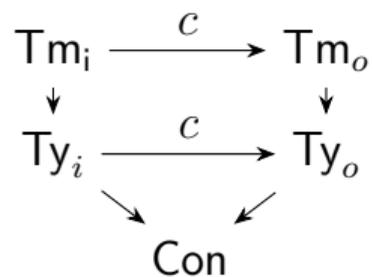
In simplicial type theory:

- ▶ Assume a bounded total order (I, \leq, \perp, \top) (or a variation)
- ▶ I is a 1-simplex (“line”), $\Delta^1 := I$
- ▶ Other simplicial shapes can be constructed, e.g. $\Delta^2 := \{(t_1, t_2) : I \times I \mid t_2 \leq t_1\}$

In 2-level type theory:

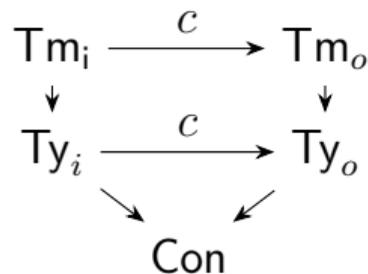
- ▶ We can mirror the STT approach and add cofibrancy assumptions.
- ▶ Alternatively:
On the outer level, define S to be the subcategory of simplicial sets, spanned by subfunctors of representables (these are the “shapes of interest”); then, assume a functor shape $: S \rightarrow \mathbf{U}^{\text{strict}}$ that sends monos to cofibrations.

Instantiating 2-level type theory



We have assumed that the inner type theory is HoTT.
We have no requirements (yet) on the outer type theory.

Instantiating 2-level type theory



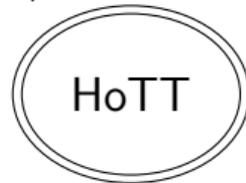
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Possibility 1: Outer theory is ext. MLTT.



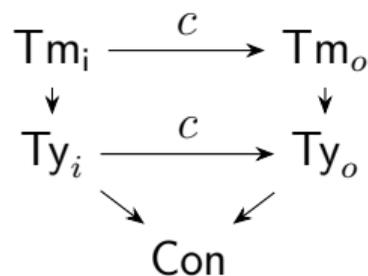
- ▶ Close to original approach [RS17]
- ▶ Less data
- ▶ Slightly more general

Possibility 2: Outer type theory is HoTT; conversion is id



- ▶ Purely in HoTT
- ▶ Matches [GWB24; GWB25]

Instantiating 2-level type theory



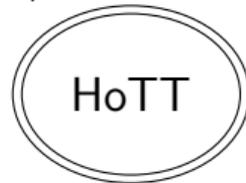
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Approaches “equivalent” b/c, for $\Sigma_X Y \rightarrow X$ and $x_0 : X$, we have:
strict fibre $\equiv Y(x_0) \simeq \Sigma(x : X).Y(x) \times (x = x_0) \equiv$ homotopy fibre

References I

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