

Arrow algebras

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Background. An *elementary topos* can be seen as a model of a version of the Calculus of Constructions with an impredicative universe of propositions, where any two elements of a proposition are definitionally equal. There is an extensive literature on topos theory (see, for example, [10, 8, 9]) and many properties of this type theory can be proved using its topos-theoretic semantics.

An important class of such toposes is the one obtained from *locales*. A locale L is a complete poset in which the following distributive law holds:

$$a \wedge \bigvee_{b \in B} b = \bigvee_{b \in B} a \wedge b,$$

when $a \in L$ and $B \subseteq L$. Every topological space gives rise to a locale by considering its poset of open subsets ordered by inclusion.

Whenever you have a locale, you can obtain a topos from it by taking the category of *sheaves* over the locale: the result is called a *localic topos*. This category of sheaves over the locale L is equivalent to a category that has a description in terms of logic. Indeed, there is an equivalent category of L -sets, which are sets with an L -valued equality relation on them, where this equality relation is required to be symmetric and transitive; the morphisms of L -sets are L -valued functional relations.

The latter category can be understood as the result of a two-step process. First, one builds a *tripos* out of the locale L and then one turns this tripos into a topos by the *tripos-to-topos construction* [6]. Importantly, there are triposes that do not arise from locales, for instance, the effective tripos, whose associated elementary topos is Hyland's effective topos, a non-localic (even non-Grothendieck) topos [5]. The effective topos and their subcategories are important as models of polymorphic type theories [4, 7].

Contribution. The aim of this talk is to introduce *arrow algebras* and explain the work of my former MSc students Marcus Briët and Umberto Tarantino [1, 15]. Arrow algebras are algebraic structures generalising locales. The point is that they still allow you to construct a tripos, an *arrow tripos*, and hence also an *arrow topos* by the tripos-to-topos construction.

These arrow toposes include the localic toposes, but also Hyland’s effective topos. Indeed, many realizability toposes can be shown to be arrow toposes, because every *pca* (*partial combinatory algebra*) gives rise to an arrow algebra: this includes also “relative, ordered” pcas as in, for example, Zoethout’s PhD thesis [16] (see also [3, 13]).

Crucially, Umberto Tarantino has developed a notion of morphism of arrow algebras which correspond to geometric morphisms between the associated triposes. This has allowed us to understand the following in purely arrow algebraic terms:

1. Every arrow morphism factors as a surjection followed by an inclusion, inducing the corresponding factorisation on the level of triposes and toposes.
2. Every subtripos of an arrow tripos coming from an arrow algebra L is induced by a *nucleus* on L . Given this nucleus, there is a simple construction of a new arrow algebra inducing the subtripos.

As a result, arrow algebras provide a flexible framework for constructing and studying new toposes.

Related work. Arrow algebras can be defined as follows:

Definition 0.1. *An arrow algebra A is a complete lattice (A, \preceq) with an implication operator $\rightarrow: A^{\text{op}} \times A \rightarrow A$ and a separator $S \subseteq A$ such that:*

1. *if $a \in S$ and $a \preceq b$, then $b \in S$.*
2. *if $a, a \rightarrow b \in S$, then also $b \in S$.*
3. *S contains the following combinators:*

$$\begin{aligned} \mathbf{k} &:= \bigwedge_{a,b} a \rightarrow b \rightarrow a \\ \mathbf{s} &:= \bigwedge_{a,b,c} (a \rightarrow b \rightarrow c) \rightarrow (a \rightarrow b) \rightarrow (a \rightarrow c) \\ \mathbf{a} &:= \bigwedge_{a, (b_i)_{i \in I}, (c_i)_{i \in I}} \left(\bigwedge_{i \in I} a \rightarrow b_i \rightarrow c_i \right) \rightarrow a \rightarrow \left(\bigwedge_{i \in I} b_i \rightarrow c_i \right) \end{aligned}$$

Arrow algebras were directly inspired by Alexandre Miquel’s work on *implicative algebras* [12]. His implicative algebras can be defined as arrow algebras in which the following axiom holds:

$$a \rightarrow \bigwedge_{b \in B} b = \bigwedge_{b \in B} a \rightarrow b.$$

We felt it was worthwhile to drop this axiom, because there are many natural examples of arrow algebras that do not satisfy it: this includes arrow algebras obtained from pcas and the arrow algebras obtained from nuclei. While it follows from Miquel’s work that every arrow algebra is equivalent to an implicative algebra [1, 11], the equivalent implicative algebra is rather unwieldy and for doing concrete calculations, working with the original arrow algebra is easier.

Implicative algebras are closely related to *evidenced frames*, as in [2]. Another framework which subsumes both realizability and localic toposes is the work by Pieter Hofstra on *BCOs*

(*basic combinatory objects*) [3] (see also [14]). While every implicative algebra is a BCO, it is not clear how to obtain a BCO from an arrow algebra in such a way that the associated triposes are isomorphic. However, a more thorough investigation of these connections is left to future work.

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